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SERIES 5

**"MATHEMATICS,
INFORMATICS AND
PHYSICS"**

VOLUME 7

CONTENTS

Mathematics

- Julia Chaparova, Eli Kalcheva* 7
Three approaches toward sublinear emden – fowler equation
- Tihomir Gyulov, Gheorghe Moroşanu* 14
On a fourth order differential inclusion involving p-biharmonic operator
- Petar Rashkov* 22
A note on averaging in differential equations with Hukuhara derivative and delay
- Tsetska Rashkova, Antoaneta Mihova* 30
Vishne identities for $M_2(G)$ and their computer realization by Mathematica

Informatics

- Desislava Atanasova* 38
An approach of visualization of belt pulley calculation results
- Galina Atanasova* 42
Conventional knowledge testing in comparison with intelligent test in algorithm area – an experimental study
- Elena Evtimova, Galina Krumova* 50
Learning development cycle applied to physics through adaptive learning methodology
- Svetlin Stoyanov, Galina Krumova* 57
Multilanguage and mobile training of physics with Moodle
- Svetlin Stoyanov, Galina Krumova, Elena Evtimova* 61
Information technologies for optimization of the learning process in the universities
- Tzvetomir Vassilev, Stanislav Kostadinov* 67
An approach to 3d on the web using Java OpenGL

Physics

- František Látal, Renata Holubová* 73
Remote experiments – new approaches to physical experimentation
- Nadezhda Nancheva, Svetlin Stoyanov* 79
M-learning of the superconductivity
- Vladimir Voinov, Roza Voinova, Zlatka Mateva* 84
An application of the extended mirror images method in constructing equivalent models for calculating the vector potential of the magnetic field produced from current fibre in inhomogeneous dielectric medium

VISHNE IDENTITIES FOR $M_2(G)$ AND THEIR COMPUTER REALIZATION BY *MATHEMATICA*

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Abstract: Vishne gave in [9] the explicit form of two identities of degree 8 for the matrix algebra $M_2(G)$, where G is the Grassmann algebra. In the paper a programme in *Mathematica* is used for proving these two identities.

Keywords: Grassmann algebra, standard polynomial, Vishne identities.

PRELIMINARIES

Let G denote the infinite dimensional Grassmann algebra, namely

$$G = G(V) = K\langle v_1, v_2, \dots \mid v_i v_j + v_j v_i = 0, i, j = 1, 2, \dots \rangle.$$

The field K has a characteristic zero. The algebra G' (without 1) has a basis $v_{i_1} v_{i_2} \dots v_{i_k}$, where $1 \leq i_1 < i_2 < \dots < i_k$. The elements v_i are called generators of G' while the elements $v_{i_1} v_{i_2} \dots v_{i_k}$ for $1 \leq i_1 < i_2 < \dots < i_k$ are called basic monomials of G' . For $G = G' \cup 1$ a generator is 1 as well. The algebras G and G' are PI-equivalent (they satisfy one and the same identities).

The algebra G is in the mainstream of recent research in PI theory. Its importance is connected with the structure theory for the T -ideals of identities of associative algebras developed by Kemer. In [4, Theorem 1.2] he proved that any T -prime T -ideal can be obtained as the T -ideal of identities of one of three algebras, one of which is the algebra $M_n(G)$.

Well known facts concerning the algebra G are the following:

Proposition 1 [5, Corollary, p. 437] *The T -ideal $T(G)$ is generated by the identity $[x_1, x_2, x_3] = 0$.*

Proposition 2 [1, Lemma 6.1] *The algebra G satisfies $S_n(x_1, \dots, x_n)^k = 0$ for all $n, k \geq 2$.*

Proposition 3 [2, Exercise 5.3] *For $G_k = G(V_k)$ over k -dimensional vector space V_k all identities follow from the identity $[x_1, x_2, x_3] = 0$ and the standard identity*

$$S_{2p}(x_1, \dots, x_{2p}) = \sum_{\sigma \in \text{Sym}(2p)} (-1)^\sigma x_{\sigma(1)} \dots x_{\sigma(2p)} = 0,$$

where p is the minimal integer such that $2p > k$.

Proposition 4 [3, Theorem 3.5] *Let K be an infinite field. A basis of the identities of G_{2k} is given by the polynomials*

$$[x_1, x_2, x_3] = 0, \quad [x_1, x_2] \dots [x_{2k+1}, x_{2k+2}] = 0.$$

Some facts concerning the identities for the matrix algebra $M_2(G)$ were considered in [7]. There it was proved that the algebra $M_n(G)$ has no identities of degree

$4n - 2$. In [8] Vishne described an efficient way to use the $Sym(n)$ -module structure of the ideal of multilinear identities in the computation of such identities of degree n of a given algebra. The method was used to show that $M_2(G)$ has identities of degree 8, but of no smaller degree. Two explicit identities of degree 8 were given. Details on the method used one could see in [8]. Here we give the definition of the identities as done in [8].

The rank and the dimension of the ideal I of 8-degree multilinear identities of $M_2(G)$ were computed and the irreducible representations of the symmetric group $Sym(8)$ were considered. There were found 15 non-zero components of I and in four cases (corresponding to the partitions $8 \vdash (2,2,1,1,1,1)$, $8 \vdash (2,2,2,1,1)$, $8 \vdash (3,1,1,1,1,1)$ and $8 \vdash (4,1,1,1,1)$) the explicit identities could be presented.

This could be done using the relationship of the $Sym(n)$ -module and the GL_m -module structures of the considered ideal for a partition $\lambda \vdash (\lambda_1, \dots, \lambda_m)$ of n [9].

Here we give only the definition of the highest weight vector of the irreducible GL_m -module. It is a non-zero element

$$f_\lambda(x_1, \dots, x_m) = \left(\prod_{i=1}^r S_{q_i}(x_1, \dots, x_{q_i}) \right) \sum_{\sigma \in Sym(n)} \alpha_\sigma \sigma$$

for some $\alpha_\sigma \in K$, where q_1, \dots, q_r are the lengths of the columns of the Young diagram related to λ .

For example for a partition $8 \vdash (2,2,1,1,1,1)$ the lengths of the columns of the corresponding Young diagram are 6 and 2. This explains the construction of the multilinear polynomials $T_1(x_1, \dots, x_6; y_1, y_2)$ and $T_2(x_1, \dots, x_5; y_1, y_2, y_3)$ done by Vishne and given in the forthcoming exposition.

A *pattern* is a finite sequence of the letter A, B . If π is a pattern with a appearances of A and b of B , we denote by $\pi(x_1, \dots, x_a; y_1, \dots, y_b)$ the product of variables where the x 's and y 's are combined according to π . For example $ABBA(x_1, x_2; y_1, y_2) = x_1 y_1 y_2 x_2$. A coefficient in front of a pattern π means that the monomial should be multiplied by that coefficient.

Now let

$$P_\pi^+ = \sum_{\sigma \in Sym(a), \tau \in Sym(b)} \text{sign}(\sigma) \pi(x_{\sigma(1)}, \dots, x_{\sigma(a)}; y_{\tau(1)}, \dots, y_{\tau(b)}),$$

$$P_\pi^- = \sum_{\sigma \in Sym(a), \tau \in Sym(b)} \text{sign}(\sigma) \text{sign}(\tau) \pi(x_{\sigma(1)}, \dots, x_{\sigma(a)}; y_{\tau(1)}, \dots, y_{\tau(b)}).$$

Let

$$P = \begin{pmatrix} + AAAABAAB, & + AABBAAAA, & - AABAAAB, \\ - AAAABBAA, & - BAABAAAA, & + BAAAABAA \end{pmatrix}.$$

The component of I in the representation $8 \vdash (2,2,1,1,1,1)$ contains (and is thus generated by) $\sum_{\pi \in P} P_\pi^-(x_1, \dots, x_6; y_1, y_2)$. Similarly, the component of I in the representation $8 \vdash (3,1,1,1,1,1)$ contains $\sum_{\pi \in P} P_\pi^+(x_1, \dots, x_6; y_1, y_2)$. In the sum

$$T_1(x_1, \dots, x_6; y_1, y_2) = \sum_{\pi \in P} (P_\pi^- + P_\pi^+) \quad (1)$$

only the monomials with y_1 preceding y_2 appear. On the other hand the original two identities are given back by $T_1(\dots; y_1, y_2) \pm T_1(\dots; y_2, y_1)$.

The same phenomenon happens for another couple of representations.

Let

$$PP = \begin{pmatrix} -AAABAABB, & -AABBAABA, & +ABBAABAA, \\ +AAABBAAB, & +AABAABBA, & -ABAABBAA \\ -ABBAAAAB, & +BAABBAAA, & -BAAAABBA \\ +ABAAAABB, & -BBAABAAA, & +BBAAAABA \end{pmatrix}.$$

The component in $8 \vdash (2,2,2,1,1)$ contains $\sum_{\pi \in PP} P_\pi^-(x_1, \dots, x_5; y_1, y_2, y_3)$ and the component in $8 \vdash (4,1,1,1,1)$ contains $\sum_{\pi \in PP} P_\pi^+(x_1, \dots, x_5; y_1, y_2, y_3)$. Again

$$T_2(x_1, \dots, x_5; y_1, y_2, y_3) = \sum_{\pi \in PP} (P_\pi^- + P_\pi^+) \quad (2)$$

has only the monomials in which the order of y_1, y_2, y_3 is even and $T_2(\dots; y_1, y_2, y_3) \pm T_2(\dots; y_3, y_2, y_1)$ gives the original identities.

Theorem 1 [8, Corollary 4.2] T_1 and T_2 are multilinear identities of degree 8 of $M_2(G)$.

COMPUTER REALIZATIONS OF THE IDENTITIES

Calculations in the Grassmann algebra are not done easily. We considered the problem of finding a computer realization of the multiplication in it. Using the 1-1 correspondence between the integer numbers from 0 to 2^n and the basic elements of a Grassmann algebra over a n-dimensional vector space a programme in *Mathematica* was written [6] using which we could prove in a computer way Vishne identities. For a guide book in the system *Mathematica* we use [9]. We point that the programme considers a finite dimensional Grassmann algebra. But this is not a limitation. Knowing the degree of the polynomial, say n, it is enough to work in the algebra G_n .

Firstly we introduce the polynomial $T_1 = T1(x1, x2, x3, x4, x5, x6, y1, y2)$ according to (1) done by *Mathematica*. We use the notation \otimes for the Grassmann multiplication.

It is easier to present the polynomial in parts corresponding to the parts of the pattern P.

Let we consider $AAAABAAB$. The corresponding part of the polynomial $T_1 = T1(x1, x2, x3, x4, x5, x6, y1, y2)$ is denoted as $A[x1, x2, x3, x4, x5, x6, y1, y2]$. Then we have:

$$S2[x1, x2] := x1 \otimes x2 - x2 \otimes x1;$$

$$S3[x1, x2, x3] := S2[x1, x2] \otimes x3 + S2[x2, x3] \otimes x1 + S2[x3, x1] \otimes x2;$$

$$S4[x1, x2, x3, x4] := S3[x1, x2, x3] \otimes x4 - S3[x2, x3, x4] \otimes x1 + S3[x3, x4, x1] \otimes x2 - S3[x4, x1, x2] \otimes x3;$$

$$A1[x1, x2, x3, x4, y1] := S4[x1, x2, x3, x4] \otimes y1;$$

$$A2[x1, x2, x3, x4, x5, y1] := A1[x1, x2, x3, x4, y1] \otimes x5 + A1[x2, x3, x4, x5, y1] \otimes x1 + A1[x3, x4, x5, x1, y1] \otimes x2 + A1[x4, x5, x1, x2, y1] \otimes x3 + A1[x5, x1, x2, x3, y1] \otimes x4;$$

$$A3[x1, x2, x3, x4, x5, x6, y1] := A2[x1, x2, x3, x4, x5, y1] \otimes x6 - A2[x2, x3, x4, x5, x6, y1] \otimes x1 + A2[x3, x4, x5, x6, x1, y1] \otimes x2 - A2[x4, x5, x6, x1, x2, y1] \otimes x3 + A2[x5, x6, x1, x2, x3, y1] \otimes x4 - A2[x6, x1, x2, x3, x4, y1] \otimes x5;$$

$$A[x1, x2, x3, x4, x5, x6, y1, y2] := A3[x1, x2, x3, x4, x5, x6, y1] \otimes y2;$$

For *AABBAAAA* the corresponding polynomial is $B[x1, x2, x3, x4, x5, x6, y1, y2]$.

We construct

$$B1[x1, x2, y1, y2] := (S2[x1, x2] \otimes y1) \otimes y2;$$

$$B2[x1, x2, x3, y1, y2] := B1[x1, x2, y1, y2] \otimes x3 + B1[x2, x3, y1, y2] \otimes x1 + B1[x3, x1, y1, y2] \otimes x2;$$

$$B3[x1, x2, x3, x4, y1, y2] := B2[x1, x2, x3, y1, y2] \otimes x4 - B2[x2, x3, x4, y1, y2] \otimes x1 + B2[x3, x4, x1, y1, y2] \otimes x2 - B2[x4, x1, x2, y1, y2] \otimes x3;$$

$$B4[x1, x2, x3, x4, x5, y1, y2] := B3[x1, x2, x3, x4, y1, y2] \otimes x5 + B3[x2, x3, x4, x5, y1, y2] \otimes x1 + B3[x3, x4, x5, x1, y1, y2] \otimes x2 + B3[x4, x5, x1, x2, y1, y2] \otimes x3 + B3[x5, x1, x2, x3, y1, y2] \otimes x4;$$

$$B[x1, x2, x3, x4, x5, x6, y1, y2] := B4[x1, x2, x3, x4, x5, y1, y2] \otimes x6 - B4[x2, x3, x4, x5, x6, y1, y2] \otimes x1 + B4[x3, x4, x5, x6, x1, y1, y2] \otimes x2 - B4[x4, x5, x6, x1, x2, y1, y2] \otimes x3 + B4[x5, x6, x1, x2, x3, y1, y2] \otimes x4 - B4[x6, x1, x2, x3, x4, y1, y2] \otimes x5;$$

For *AABAAAB* the corresponding polynomial is $CH[x1, x2, x3, x4, x5, x6, y1, y2]$.

We have

$$\begin{aligned} CH1[x1, x2, y1] &:= S2[x1, x2] \otimes y1; \\ CH2[x1, x2, x3, y1] &:= CH1[x1, x2, y1] \otimes x3 + CH1[x2, x3, y1] \otimes x1 + \\ &CH1[x3, x1, y1] \otimes x2; \end{aligned}$$

$$\begin{aligned} CH3[x1, x2, x3, x4, y1] &:= CH2[x1, x2, x3, y1] \otimes x4 - \\ &CH2[x2, x3, x4, y1] \otimes x1 + CH2[x3, x4, x1, y1] \otimes x2 - \\ &CH2[x4, x1, x2, y1] \otimes x3; \end{aligned}$$

$$\begin{aligned} CH4[x1, x2, x3, x4, x5, y1] &:= CH3[x1, x2, x3, x4, y1] \otimes x5 + \\ &CH3[x2, x3, x4, x5, y1] \otimes x1 + CH3[x3, x4, x5, x1, y1] \otimes x2 + \\ &CH3[x4, x5, x1, x2, y1] \otimes x3 + CH3[x5, x1, x2, x3, y1] \otimes x4; \end{aligned}$$

$$\begin{aligned} CH5[x1, x2, x3, x4, x5, x6, y1] &:= CH4[x1, x2, x3, x4, x5, y1] \otimes x6 - \\ &CH4[x2, x3, x4, x5, x6, y1] \otimes x1 + CH4[x3, x4, x5, x6, x1, y1] \otimes x2 - \\ &CH4[x4, x5, x6, x1, x2, y1] \otimes x3 + CH4[x5, x6, x1, x2, x3, y1] \otimes x4 - \\ &CH4[x6, x1, x2, x3, x4, y1] \otimes x5; \end{aligned}$$

$$CH[x1, x2, x3, x4, x5, x6, y1, y2] := CH5[x1, x2, x3, x4, x5, x6, y1] \otimes y2;$$

The corresponding part to *AAAABBAA* is denoted as $DH[x1, x2, x3, x4, x5, x6, y1, y2]$. We get

$$DH1[x1, x2, x3, x4, y1, y2] := (S4[x1, x2, x3, x4] \otimes y1) \otimes y2;$$

$$\begin{aligned} DH2[x1, x2, x3, x4, x5, y1, y2] &:= DH1[x1, x2, x3, x4, y1, y2] \otimes x5 + \\ &DH1[x2, x3, x4, x5, y1, y2] \otimes x1 + DH1[x3, x4, x5, x1, y1, y2] \otimes x2 + \\ &DH1[x4, x5, x1, x2, y1, y2] \otimes x3 + DH1[x5, x1, x2, x3, y1, y2] \otimes x4; \end{aligned}$$

$$\begin{aligned} DH[x1, x2, x3, x4, x5, x6, y1, y2] &:= DH2[x1, x2, x3, x4, x5, y1, y2] \otimes x6 - \\ &DH2[x2, x3, x4, x5, x6, y1, y2] \otimes x1 + DH2[x3, x4, x5, x6, x1, y1, y2] \otimes x2 - \\ &DH2[x4, x5, x6, x1, x2, y1, y2] \otimes x3 + DH2[x5, x6, x1, x2, x3, y1, y2] \otimes x4 - \\ &DH2[x6, x1, x2, x3, x4, y1, y2] \otimes x5; \end{aligned}$$

The part *BAABAAAA* gives rise to $EH[x1, x2, x3, x4, x5, x6, y1, y2]$. Thus

$$\begin{aligned} EH1[x1, x2, y1, y2] &:= ((y1 \otimes x1) \otimes x2) \otimes y2 - ((y1 \otimes x2) \otimes x1) \otimes y2; \\ EH2[x1, x2, x3, y1, y2] &:= EH1[x1, x2, y1, y2] \otimes x3 + \\ &EH1[x2, x3, y1, y2] \otimes x1 + EH1[x3, x1, y1, y2] \otimes x2; \end{aligned}$$

$$\begin{aligned} EH3[x1, x2, x3, x4, y1, y2] &:= EH2[x1, x2, x3, y1, y2] \otimes x4 - \\ &EH2[x2, x3, x4, y1, y2] \otimes x1 + EH2[x3, x4, x1, y1, y2] \otimes x2 - \\ &EH2[x4, x1, x2, y1, y2] \otimes x3; \end{aligned}$$

$$\begin{aligned} EH4[x1, x2, x3, x4, x5, y1, y2] &:= EH3[x1, x2, x3, x4, y1, y2] \otimes x5 + \\ &EH3[x2, x3, x4, x5, y1, y2] \otimes x1 + EH3[x3, x4, x5, x1, y1, y2] \otimes x2 + \\ &EH3[x4, x5, x1, x2, y1, y2] \otimes x3 + EH3[x5, x1, x2, x3, y1, y2] \otimes x4; \end{aligned}$$

$$\begin{aligned} EH[x1, x2, x3, x4, x5, x6, y1, y2] &:= EH4[x1, x2, x3, x4, x5, y1, y2] \otimes x6 - \\ &EH4[x2, x3, x4, x5, x6, y1, y2] \otimes x1 + EH4[x3, x4, x5, x6, x1, y1, y2] \otimes x2 - \\ &EH4[x4, x5, x6, x1, x2, y1, y2] \otimes x3 + EH4[x5, x6, x1, x2, x3, y1, y2] \otimes x4 - \\ &EH4[x6, x1, x2, x3, x4, y1, y2] \otimes x5; \end{aligned}$$

The polynomial $FH[x1, x2, x3, x4, x5, x6, y1, y2]$ is related to $BAAAABAA$ and is constructed in a similar way, namely

$$FH1[x1, x2, y1] := (y1 \otimes x1) \otimes x2 - (y1 \otimes x2) \otimes x1;$$

$$\begin{aligned} FH2[x1, x2, x3, y1] &:= FH1[x1, x2, y1] \otimes x3 + FH1[x2, x3, y1] \otimes x1 + \\ &FH1[x3, x1, y1] \otimes x2; \end{aligned}$$

$$\begin{aligned} FH3[x1, x2, x3, x4, y1] &:= FH2[x1, x2, x3, y1] \otimes x4 - \\ &FH2[x2, x3, x4, y1] \otimes x1 + FH2[x3, x4, x1, y1] \otimes x2 - \\ &FH2[x4, x1, x2, y1] \otimes x3; \end{aligned}$$

$$FH4[x1, x2, x3, x4, y1, y2] := FH3[x1, x2, x3, x4, y1] \otimes y2;$$

$$\begin{aligned} FH5[x1, x2, x3, x4, x5, y1, y2] &:= FH4[x1, x2, x3, x4, y1, y2] \otimes x5 + \\ &FH4[x2, x3, x4, x5, y1, y2] \otimes x1 + FH4[x3, x4, x5, x1, y1, y2] \otimes x2 + \\ &FH4[x4, x5, x1, x2, y1, y2] \otimes x3 + FH4[x5, x1, x2, x3, y1, y2] \otimes x4; \end{aligned}$$

$$\begin{aligned} FH[x1, x2, x3, x4, x5, x6, y1, y2] &:= FH5[x1, x2, x3, x4, x5, y1, y2] \otimes x6 - \\ &FH5[x2, x3, x4, x5, x6, y1, y2] \otimes x1 + FH5[x3, x4, x5, x6, x1, y1, y2] \otimes x2 - \\ &FH5[x4, x5, x6, x1, x2, y1, y2] \otimes x3 + FH5[x5, x6, x1, x2, x3, y1, y2] \otimes x4 - \\ &FH5[x6, x1, x2, x3, x4, y1, y2] \otimes x5; \end{aligned}$$

At the end we form the entire polynomial

$$T1[x1, x2, x3, x4, x5, x6, y1, y2] := A[x1, x2, x3, x4, x5, x6, y1, y2] + \\ B[x1, x2, x3, x4, x5, x6, y1, y2] - CH[x1, x2, x3, x4, x5, x6, y1, y2] - \\ DH[x1, x2, x3, x4, x5, x6, y1, y2] - EH[x1, x2, x3, x4, x5, x6, y1, y2] + \\ FH[x1, x2, x3, x4, x5, x6, y1, y2];$$

In a similar way the recurrent construction of $T_2 = T2(x1, x2, x3, x4, x5, y1, y2, y3)$ is realized due to (2).

Then we set $n = 8$ in the *Mathematica* programme realized in [6]. The evaluation on $M_2(G)$ was done in two stages.

We use the operator **Random[type, {min, max}]** bringing out an arbitrary number of type Integer, Real or Complex. At the beginning we specify the type as **Real**. Due to the rounding done in calculations the needed result was not obtained.

Then for exact calculations we specify the type as **Integer**. For max=500 for example we have

$$\text{For}[i = 1, i \leq 2^8, i_{++}, \{a[i] = \text{Random}[\text{Integer}, \{0, 500\}], b[i] = \text{Random}[\text{Integer}, \{0, 500\}], \\ c[i] = \text{Random}[\text{Integer}, \{0, 500\}], d[i] = \text{Random}[\text{Integer}, \{0, 500\}]\}]$$

for a 2×2 matrix

$$x = \{ \{ \text{Array}[a, 256], \text{Array}[b, 256] \}, \{ \text{Array}[c, 256], \text{Array}[d, 256] \} \}.$$

For the evaluation of the polynomial $T_1(x_1, \dots, x_6, y_1, y_2)$ by Intel Pentium computer with 2GB RAM 40 minutes were needed.

The calculation of $T_2(x_1, \dots, x_5, y_1, y_2, y_3)$ with random matrices took 70 minutes time.

Then we consider the general case with arbitrary matrices introducing a matrix variable x only as

$$x = \{ \{ \text{Array}[a, 256], \text{Array}[b, 256] \}, \{ \text{Array}[c, 256], \text{Array}[d, 256] \} \}.$$

The possibilities of our Pentium computer were not enough for calculating the polynomials $T_1(x_1, \dots, x_6, y_1, y_2)$ and $T_2(x_1, \dots, x_5, y_1, y_2, y_3)$ in the general case.

We point that the identity $[x_1, x_2, x_3]^2 = 0$ for the upper triangular two by two matrices with entries from G_n for $n = 12$ was confirmed for 3 hours by Intel Celeron computer with 2GB RAM.

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ТЪЖДЕСТВА НА ВИШНЕ ЗА $M_2(G)$ И ТЯХНАТА КОМПЮТЪРНА РЕАЛИЗАЦИЯ ЧРЕЗ *MATHEMATICA*

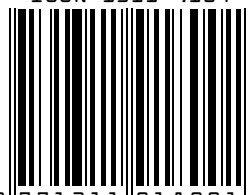
Цецка Рашкова, Антоанета Михова

Русенски университет "Ангел Кънчев"

Резюме: В своя работа [8] Вишне дава явната форма на две тъждества от степен 8 за матричната алгебра $M_2(G)$, където G е Грасмановата алгебра. В статията се използва програма на *Mathematica* за доказване на тези две тъждества.

Ключови думи: Грасманова алгебра, стандартен полином, тъждества на Вишне

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